

LANCOM Techpaper

Routing Performance

Applications for communications and entertainment are increasingly based on IP networks. In order to ensure that the necessary bandwidth performance can be provided reliably, it is important for the infrastructure's networking components to be tested thoroughly and intensively. In this techpaper, LANCOM Systems presents the methods of measuring routing and VPN performance for central site and VPN gateways as well as the respective results.

We have examined a variety of aspects for consideration when measuring the router performance. This includes transmission speeds of connections between the LAN and the Internet (WAN), and the internal data transmission in the network (LAN-LAN). Many business processes rely on secure WAN connections, which is why we have focused on determining the performance of encrypted data connections over VPN.

Test System

All of the performance values were measured in the LANCOM test laboratory. Tests were conducted with an IXIA test system. IXIA uses so-called test suites, which enable the simulation of different applications. This allows, for example, the investigation of data throughput over automatically established VPN tunnels, or the measurement of pure LAN-WAN routing performance for unidirectional and bidirectional data connections. IXIA is a leading supplier of systems which test IP-based services and infrastructures. Test systems from IXIA are employed all over the world by network component manufacturers and other organizations to help assure the functionality and reliability of complex IP networks, devices, and applications.

The measurement of data transmission itself uses either a fixed frame size or a combination of frame sizes which reflects a typical flow of data. These combinations are known as "Internet Mix", or IMIX for short. The type of

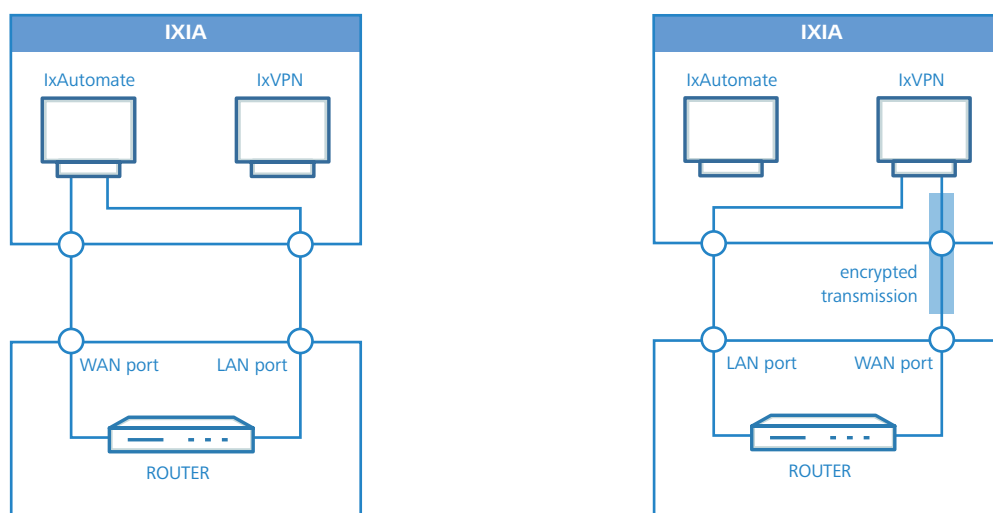


Figure 1: IXIA test system for routing connections and encrypted VPN connections between LAN and WAN

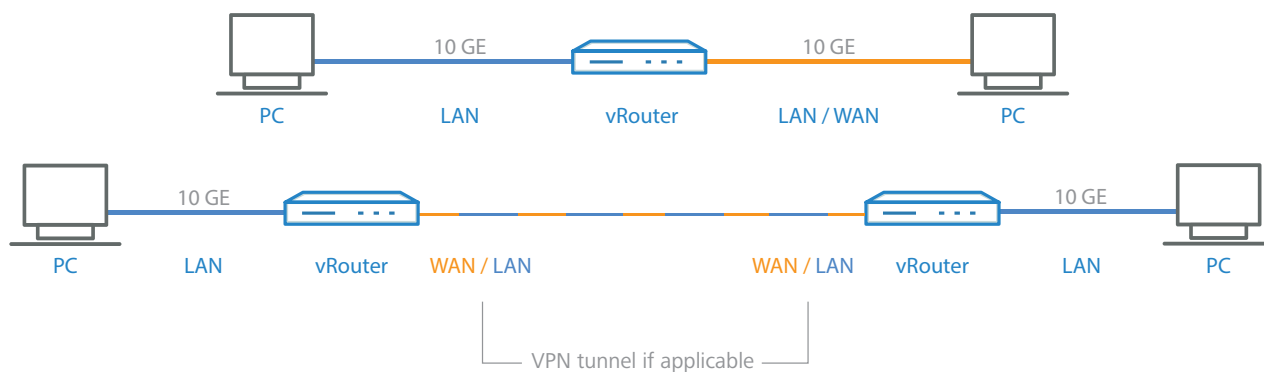


Figure 2: Schematic view of the vRouter test system scenarios

IMIX which is applied significantly affects the test results because packet size has a strong influence on a connection’s performance. By selecting the appropriate ports on the router being tested, it is possible to test connections between the LAN and the WAN, and also pure LAN-LAN connections.

Two scenarios are used for the vRouter measurements. One is the transmission from a PC in the LAN to a PC in another LAN or WAN by using the vRouter in between. The second scenario adds two vRouters which are connected over WAN and use a VPN tunnel for encrypted data transfer.

Firmware on the devices under test is LCOS 10.20 or LCOS 8.70 CC for the LANCOM CC VPN router. vRouter measurements were done with LCOS 10.12.

Routing Performance (UDP)

The measurement of routing performance involves the determination of the maximum data throughput which can be achieved before a router starts rejecting packets. This measurement uses UDP packets of various sizes in order to simulate the performance with different applications. Ethernet frame sizes range from 64 bytes for the smallest to 1518 bytes for the largest frames. Tests on different router models demonstrate the influence of the different hardware platforms (processor, interfaces).

Measurements initially determine the frame rate, which

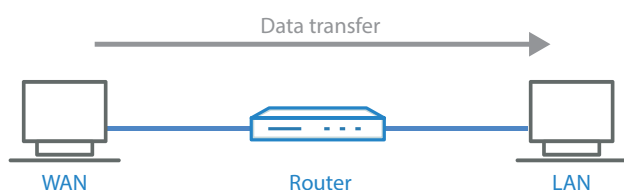


Figure 3: Schematic view of the test system

is a good performance indicator of the tested hardware. With normal routing, the frame rate is fairly constant even with different frame sizes. This is because only the header is inspected during routing, a process which is largely independent of the size of the frames being routed. For this reason, only the average frame rates are given in the tables.

The throughput for a certain frame size (or even a mix of sizes, see IMIX on page 4) can be approximately calculated by multiplication with the frame rate. When the frame rate is constant, data throughput depends directly on the frame size because the larger the frames, the larger is the data volume that can be transmitted. The maximum number of frames transmitted per second is limited by the performance of the interfaces and the transmission medium.

Measurement of the routing performance relates to the size of the Ethernet frames. To compare packet sizes for particular applications, it is necessary to subtract the header. For a frame of 512 bytes, the result is a UDP datagram size of 474 bytes (512 bytes - 14 bytes Ethernet header - 4 bytes FCS trailer - 20 bytes IP header) and, after subtracting the UDP header (8 bytes), the UDP payload is 466 bytes.

To investigate routing performance, in this paper two different applications are considered.

- > For WAN-LAN routing, data received from the WAN is forwarded to a peer in the LAN.
- > For LAN-LAN routing, data remains within the local-area network and is passed from one LAN port to another.

The measurements show that the throughput increases almost linearly with the frame size until the limit of the Gigabit interface is reached.

Routing Performance (TCP)

Using UDP, the maximum performance is easily measured. However, a large portion of data transfers utilizes TCP. Therefore, it is important to take a closer look at those scenarios as well.

The TCP measurement is accomplished by using iperf, a program used to measure the data throughput between two computers. They are connected to a LANCOM router. One computer acts as a server in the WAN, sending data packets to the other, which acts as a client in the LAN, representing a typical download scenario.

The TCP measurement done with iperf did not use NAT since no direct measurement with NAT is possible using iperf. The correlation between the test results and the results seen at a WAN connection including NAT in a live environment is very high. The causes are the differences of the traffic structure of iperf compared to an HTTP download and the LCOS 10.12 firewall which offers a similar performance with NAT or without.

Both computers used in the test system have a similar hardware and software:

- > Intel Core i7 CPU
- > Intel PRO/1000 nic
- > Ubuntu 12.04 / Kernel 3.8.0

The measurement was run using iperf 2.05. The TCP window size was set to 256 kb and five sessions were run simultaneously.

vRouter measurements were done with the following hardware:

- > PCs: Core i7-6700, Intel X540 10GE interfaces
- > ESXi Server: Dell PowerEdge R330 with Xeon E3-1230v5, 3.4 GHz, Intel X710 10GE interfaces as uplink to the ESXi vSwitches
- > vRouter: VMXNET3 with virtual interfaces

IPerf was used without window parameter for the vRouter tests.

IPSec Routing Performance

Other than with pure routing performance, IPSec-VPN routing actually changes the frames which are being passed from one interface to the next. When data is encrypted for the VPN tunnel, the original frame is encapsulated and it is supplemented with additional information. This has two important effects when considering the performance of IPSec routing:

- > Encrypted frames are larger than unencrypted frames. Consequently, any results have to indicate which frame size was observed at which interface, and/or whether frames were encrypted or unencrypted. The values presented here always relate to an unencrypted frame size. An IP packet of 46 bytes is transported unencrypted, e.g. in a frame of 64 bytes. In the event of AES encryption, the frame grows for example to 122 bytes (46 byte IP packet + 14 byte Ethernet + 4 byte FCS + 20 byte IP + 8 byte ESP + 16 byte initialization vector (IV) + 0 byte padding (0-15 byte) + 1 byte padding length + 1 byte next header + 12 byte authentication).
- > The processes of encryption and decryption in the router take up computing time. These processes take place in two steps which, in the case of encryption, must be sequential. With decryption, on the other hand, these steps can be executed in parallel. Router models with VPN hardware acceleration provide significantly better performance with decryption than with encryption. This explains why the results display a significant difference in performance between the decryption and encryption directions. All of the IPSec-routing values given here are for a single VPN tunnel. With up to 1,000 tunnels established under laboratory conditions, the frame rate remained almost constant over all of the active tunnels. However, under actual operating conditions an increasing number of tunnels will cause the frame rate to drop due to the processes running for each tunnel (for example renewal of the key being used).

IPSec Routing with Different IMIXs (Decryption and Encryption)

As an alternative to measurement with fixed frame sizes, measurements were carried out with varying IMIX patterns. The IMIX patterns simulate “true” data traffic, which contains varying frame sizes. There are no fixed guidelines for the composition of the frame sizes, and for this reason the measurements supplemented the preset values in the IXIA test system (IMIX 0) with two other common patterns (IMIX 1 and IMIX 2). The patterns use the following frame compositions:

- > IMIX 0: 45% 64 byte, 20% 128 byte, 5% 256 byte, 3% 512 byte, 2% 1024 byte, 1% 1280 byte, 24% 1364 byte.
- > IMIX 1: 7x 64 byte, 4x 570 byte, 1x 1418 byte.
- > IMIX 2: 58% 90 byte, 2% 92 byte, 24% 594 byte, 16% 1418 byte.

By assuming 100 byte overhead the maximum encrypted frame size transmittable via Ethernet is 1418 bytes (1518 bytes is the maximum IEEE 802.3 frame size).

AES128-SHA encryption was used for the measurements, and the tunnels were established in the LAN-WAN direction. These measurements again show that most commonly data decryption is faster than encryption.

Bridging Performance (L2TPv3 Ethernet tunnel)

Bridging performance (Table 10) shows, which maximum data throughput is possible via a transparent Ethernet tunnel, which is established over L2TPv3 pseudowire. A measurement with 5 parallel TCP streams using IPerf is performed here through the Ethernet tunnel.

vRouter measurements were done with the exact hardware mentioned under „Routing Performance (TCP)“, using a WAN-WAN connection. The rest of the measurements is done via LAN-LAN connection.

Table 1:

WAN-LAN routing – Throughput [Mbps] per frame size [byte] and average frame rate [fps]

Device	LCOS	64	128	256	512	1024	1280	1518	frame rate
1640E									
1780EW-4G+									
1781EF+									
1781EW+									
1790-4G									
1781VA									
1783VA									
1784VA									
1790VA	10.20	47.2	94.7	188.8	376.5	716.1	893.9	984.4	89215
1793VA									
1781VAW									
1783VAW									
1790VAW									
1793VAW									
1781VA-4G									
1783VA-4G									
1790VA-4G									
1793VA-4G									
1900EF									
1906VA	10.20	73.2	147.1	294.3	585.8	977.1	981.6	984.4	124228
1906VA-4G									
ISG-1000	10.20	68.1	135.9	272.3	543.5	980.8	984.6*	987.0*	118359
WLC-1000									
ISG-4000	10.20	100.0	200.0	401.3	800.0	980.8	984.6*	987.0*	154146
1781EF (CC)	8.70 CC	5.2	10.6	21.3	42.8	85.4	106.4	125.4	10362

*Test setup limited to 1 Gbps

Table 2:

LAN-LAN routing – Throughput [Mbps] per frame size [byte] and average frame rate [fps]

Device	LCOS	64	128	256	512	1024	1280	1518	frame rate
1640E									
1780EW-4G+									
1781EF+									
1781EW+									
1790-4G									
1781VA									
1783VA									
1784VA									
1790VA	10.20	53.4	103.1	212.6	419.0	797.5	980.1	984.4	97889
1793VA									
1781VAW									
1783VAW									
1790VAW									
1793VAW									
1781VA-4G									
1783VA-4G									
1790VA-4G									
1793VA-4G									
1900EF									
1906VA	10.20	79.0	157.6	316.0	632.1	977.1	981.6	984.4	130443
1906VA-4G									
ISG-1000	10.20	82.9	165.8	331.6	664.9	980.8	984.6*	987.0*	135036
WLC-1000									
ISG-4000	10.20	113.1	226.9	453.9	904.6	980.8	984.6*	987.0*	168874
1781EF (CC)	8.70 CC	6.7	13.5	27.2	53.2	101.7	126.3	149.1	12804

*Test setup limited to 1 Gbps

Table 3:

WAN-LAN routing – Throughput [Mbps]

Device	LCOS	Throughput 5 TCP sessions
1640E		
1780EW-4G+		
1781EF+		
1781EW+		
1790-4G		
1781VA		
1783VA		
1784VA		
1790VA	10.20	739.0
1793VA		
1781VAW		
1783VAW		
1790VAW		
1793VAW		
1781VA-4G		
1783VA-4G		
1790VA-4G		
1793VA-4G		
1900EF		
1906VA	10.20	908.0
1906VA-4G		
1781EF (CC)	8.70 CC	91.7

Table 4:

IPSec routing (decryption) – Throughput [Mbps] per frame size [byte] and average frame rate [fps]

Device	LCOS	64	128	256	512	1024	1280	1418	frame rate
1640E									
1780EW-4G+									
1781EF+									
1781EW+									
1790-4G									
1781VA									
1783VA									
1784VA									
1790VA	10.20	24.6	48.5	96.9	189.8	370.9	453.4	499.3	52939
1793VA									
1781VAW									
1783VAW									
1790VAW									
1793VAW									
1781VA-4G									
1783VA-4G									
1790VA-4G									
1793VA-4G									
1900EF									
1906VA	10.20	37.2	73.9	150.2	299.2	592.2	725.4	793.8	72058
1906VA-4G									
ISG-1000	10.20	36.5	72.8	146.0	291.5	580.6	721.1	800.4	70942
WLC-1000									
ISG-4000	10.20	50.4	99.5	190.9	385.4	760.1	914.1	921.2	95295
1781EF (CC)	8.70 CC	4.1	9.1	17.5	34.9	67.3	83.6	91.4	8340

Table 5:

IPSec routing (encryption) – Throughput [Mbps] per frame size [byte] and average frame rate [fps]

Device	LCOS	64	128	256	512	1024	1280	1418	frame rate
1640E									
1780EW-4G+									
1781EF+									
1781EW+									
1790-4G									
1781VA									
1783VA									
1784VA									
1790VA	10.20	22.1	44.7	88.3	175.9	343.4	424.7	466.2	42473
1793VA									
1781VAW									
1783VAW									
1790VAW									
1793VAW									
1781VA-4G									
1783VA-4G									
1790VA-4G									
1793VA-4G									
1900EF									
1906VA	10.20	35.2	70.1	140.0	279.6	556.1	691.8	780.1	68284
1906VA-4G									
ISG-1000	10.20	33.4	66.8	134.6	268.2	535.4	666.8	741.5	65348
WLC-1000									
ISG-4000	10.20	45.9	93.6	186.5	374.4	747.1	933.3	947.6	89898
1781EF (CC)	8.70 CC	3.0	6.2	11.5	23.8	46.6	58.5	65.0	5778

*Test setup limited to 1 Gbps

Table 6:

IPSec routing (decryption) – Throughput [Mbps]

Device	LCOS	IMIX 0	IMIX 1	IMIX 2
1640E				
1780EW-4G+				
1781EF+				
1781EW+				
1790-4G				
1781VA				
1783VA				
1784VA				
1790VA	10.20	153.2	120.9	123.9
1793VA				
1781VAW				
1783VAW				
1790VAW				
1793VAW				
1781VA-4G				
1783VA-4G				
1790VA-4G				
1793VA-4G				
1900EF				
1906VA	10.20	235.1	186.3	225.4
1906VA-4G				
ISG-1000	10.20	224.4	175.6	213.7
WLC-1000				
ISG-4000	10.20	316.1	247.8	300.5
1781EF (CC)	8.70 CC	26.4	20.5	24.4

Table 7:

IPSec routing (encryption) – Throughput [Mbps]

Device	LCOS	IMIX 0	IMIX 1	IMIX 2
1640E				
1780EW-4G+				
1781EF+				
1781EW+				
1790-4G				
1781VA				
1783VA				
1784VA				
1790VA	10.20	144.3	113.1	137.5
1793VA				
1781VAW				
1783VAW				
1790VAW				
1793VAW				
1781VA-4G				
1783VA-4G				
1790VA-4G				
1793VA-4G				
1900EF				
1906VA	10.20	232.1	181.4	220.3
1906VA-4G				
ISG-1000	10.20	210.6	164.8	200.9
WLC-1000				
ISG-4000	10.20	300.4	233.1	287.7
1781EF (CC)	8.70 CC	21.1	17.2	21.4

Table 8:**TCP throughput [Mbps] (see page 2, fig. 2 top)**

LANCOM	Method	LCOS	Throughput 5 TCP sessions
vRouter	LAN - WAN (IPoE) routing	10.20	9376
vRouter	LAN - LAN routing	10.20	9373
ISG-1000	LAN - WAN (IPoE) Routing	10.20	940*
ISG-1000	LAN - LAN Routing	10.20	940*
ISG-4000	LAN - WAN (IPoE) Routing	10.20	1935
ISG-4000	LAN - LAN Routing	10.20	3345

*Limited by 1 Gbps ports of the ISG-1000

Table 9:**TCP throughput [Mbps] with VPN tunnel if applicable - (see page 2, fig.2 bottom)**

LANCOM	Method	LCOS	Throughput 5 TCP sessions
vRouter	LAN - WAN (IPoE) - LAN routing	10.20	7921
vRouter	LAN - WAN (IPSec AES256 SHA256) - LAN routing	10.20	845
vRouter	LAN - WAN (IPSec AES256 GCM) - LAN routing	10.20	2316
ISG-4000	LAN - LAN (IPoE) - LAN routing	10.20	1803
ISG-4000	LAN - LAN (IPSec AES256 SHA256) - LAN routing	10.20	1015
ISG-4000	LAN - LAN (IPSec AES256 GCM) - LAN routing	10.20	1025

Table 10:**L2TPv3 bridging - Throughput [Mbps]**

LANCOM	LCOS	Throughput 5 TCP sessions
ISG-1000	10.20	991
WLC-1000		
ISG-4000	10.20	1562
vRouter	10.20	5270